Transport Layer (II)

Error-control Method (cont.)

- The error-control method used in the transport layer is **Automatic Repeat reQuest (ARQ)**. This method has six components:
  1. **acknowledgement**
     - receiver explicitly tells sender that a packet is received OK
     - so, sender will know whether the packet is delivered or whether the packet needs to be retransmitted.
  2. **timeout**
     - how long should the sender wait for ACK? what if ACK/data is lost? how does sender know whether the packet is received or not?
     - timeout is used to solve the problem. Sender triggers a timer immediately after transmit a packet. If sender does not receive ACK within a certain time (we call it timeout), retransmit the packet.
  3. **retransmission**
     - sender retransmit the current packet if timeout
  4. **give up after k retransmissions**
     - what if the link is very bad? very lossy? [sender gives up after k retransmissions]
     - prevent infinite retransmissions

- The above components work well if DATA is lost, shown in the left figure below. What if DATA is received but ACK is lost? The sender will retransmit due to timeout. Then there will be two copies at the receiver side, which is confusing and may lead to bad user experience. The duplicate packet transmission also wastes computing and network resource. How do we solve it?
5. sequence numbers on data packets
   - handle duplicates happened in retransmission
   - sender adds a unique sequence number to each packet
   - receiver can detect duplicate packets based on their sequence number and discard
duplicate packets (shown as the right figure above)

   • So far so good. However, it's possible that ACK will be received by the sender out of order.
   If so, it's hard to tell which packets have been received. As shown in the figure below, it's
   hard to tell wether the last received ACK is for packet 7 or packet 8, and packet 8 is lost.

6. cumulative acknowledgment numbers
   - add sequence numbers on ACKs
- the sequence numbers on ACKs are not the received packet numbers but the next expected packet. This improves the efficiency, as an ACK indicates all prior packets were received even if ACK for a prior received packet is lost, show as the figure below.

- All the scenarios discussed above use the same way to fill the communication pipe, Stop and Wait.
  - In this type of protocols, senders wait for acknowledgement before seeding next packet.
- However, Sliding Window protocols are widely used in nowadays.
  - It lets senders send up to W (window size) packets before waiting for acknowledgement.

Disclaimer: Notes adapted from the textbook and online resources.